

# XMC1000

Microcontroller Series  
for Industrial Applications

Interrupt Subsystem

✓ Using Interrupts

Device Guide

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**Revision History**

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# Using Interrupts

## 1 Using Interrupts

An interrupt is an exception signaled by a peripheral, or generated by a software request.

Setting up an interrupt in a XMC1000 system involves the following steps:

- Configure the Interrupt Priority Level (optional)
- Define an Interrupt Service Routine (ISR)
- Enable the Interrupt

### 1.1 Configure the Interrupt Priority Level

This step is optional.

An interrupt node can be configured to one of four priority levels. The levels are in steps of 64, from 0 to 192, with level 0 having the highest priority. By default, all nodes are configured to level 0.

To change the interrupt priority level, use the CMSIS function `NVIC_SetPriority()`. For example, to configure node 1, which is connected to SCU service request output 1 (SCU\_SR1), to level 192, the following code can be used:

```
//This CMSIS function configures node 1 to priority level 192
NVIC_SetPriority(SCU_1_IRQn, 0xC0);
```

SCU\_1\_IRQn is the enumerated constant for node 1. The complete list of enumerated constants can be found in the device header file "XMC1x00.h".

### 1.2 Define an Interrupt Service Routine (ISR)

An interrupt node may contain more than one interrupt trigger source. For example, node 1 (SCU\_SR1) triggers an interrupt request whenever there is a standby clock failure or VDDP pre-warning event, assuming these events are enabled for interrupt generation.

When defining an ISR or interrupt handler, the following should be considered:

- Identify the handler name from the list given in the device's startup file "startup\_XMC1x00.s".
- Check the status flag of the highest priority event first to determine if it is the interrupt trigger source, and the processing task executed, before proceeding to check for the next highest priority event.

Note: The startup file handles the interrupt vector address and the method of pointing to the actual interrupt handler location.

The ISR for node 1 could take the following form:

```
void SCU_1_IRQHandler(void)
{
    //Check if standby clock failure event status flag is set
    if(0U != (SCU_INTERRUPT->SRRW & SCU_INTERRUPT_SRRW_SBYCLKFI_Msk))
    {
        //Write to SRCLR to clear active flag
        SCU_INTERRUPT->SRCLR |= SCU_INTERRUPT_SRCLR_SBYCLKFI_Msk;
        ... //Process the interrupt
    }
    //Check if VDDP pre-warning event status flag is set
    if (0U != (SCU_INTERRUPT->SRRW & SCU_INTERRUPT_SRRW_VDDPI_Msk))
    {
        //Write to SRCLR to clear active flag
```

```

        SCU_INTERRUPT->SRCLR |= SCU_INTERRUPT_SRCLR_VDDPI_Msk;
        ... //Process the interrupt
    }
    return;
}

```

For events that are not enabled for interrupt generation, there is no need to check their status flags within the interrupt handler.

Interrupt handlers are by default allocated to the Flash memory by the Integrated Development Environment (IDE) tools, such as Infineon DAVE™ 3 and µKeil Vision, during a software build.

If an interrupt service request is generated while the Flash is busy, and the CPU attempts to vector to the handler location in the Flash, a system bus stall will occur. CPU execution of the interrupt handler will only occur when the Flash is available again.

To avoid the above scenario, interrupt handlers can be relocated to the SRAM. For information on relocating code to SRAM, refer to the respective IDE tool documentation.

### 1.3 Enable the Interrupt

To generate an interrupt request to CPU, both the interrupt node and the underlying interrupt trigger sources must be enabled for interrupt request generation.

Our recommendation is to first clear the interrupt node 'pending' status and the event status flags of the interrupt trigger sources, before enabling them. Use the following steps:

- Configure the interrupt node pointer (optional)
- Clear the interrupt node pending status
- Enable the interrupt node
- Clear the event status flags
- Enable the interrupt trigger sources

The following example code sequence enables interrupt node 1 (SCU\_SR1), and the standby clock failure and VDDP pre-warning events for interrupt request generation:

```

//This function clears node 1 pending status
NVIC_ClearPendingIRQ(SCU_1_IRQn);
//This function enables node 1 for interrupt request generation
NVIC_EnableIRQ(SCU_1_IRQn);
//Initializes event status flags SBYCLKFI and VDDPI to 0
SCU_INTERRUPT->SRCLR |= SCU_INTERRUPT_SRCLR_SBYCLKFI_Msk;
SCU_INTERRUPT->SRCLR |= SCU_INTERRUPT_SRCLR_VDDPI_Msk;
//Enables SBYCLKFI and VDDPI interrupt trigger sources
SCU_INTERRUPT->SRMSK |= SCU_INTERRUPT_SRMSK_SBYCLKFI_Msk;
SCU_INTERRUPT->SRMSK |= SCU_INTERRUPT_SRMSK_VDDPI_Msk;

```

### 1.4 Processing an Interrupt with Multiple Trigger Sources

Assume that an interrupt node is assigned to the service request output of Module n (MOD\_n\_SR) and there are two events, A and B, that can generate an interrupt request. The pseudo code of the interrupt handler is given below:

```

void MOD_n_IRQHandler(void)
{
    //Check if event A status flag is set

```

```

{
    //Process the interrupt
}
//Check if event B status flag is set
{
    //Process the interrupt
}
return;
}

```

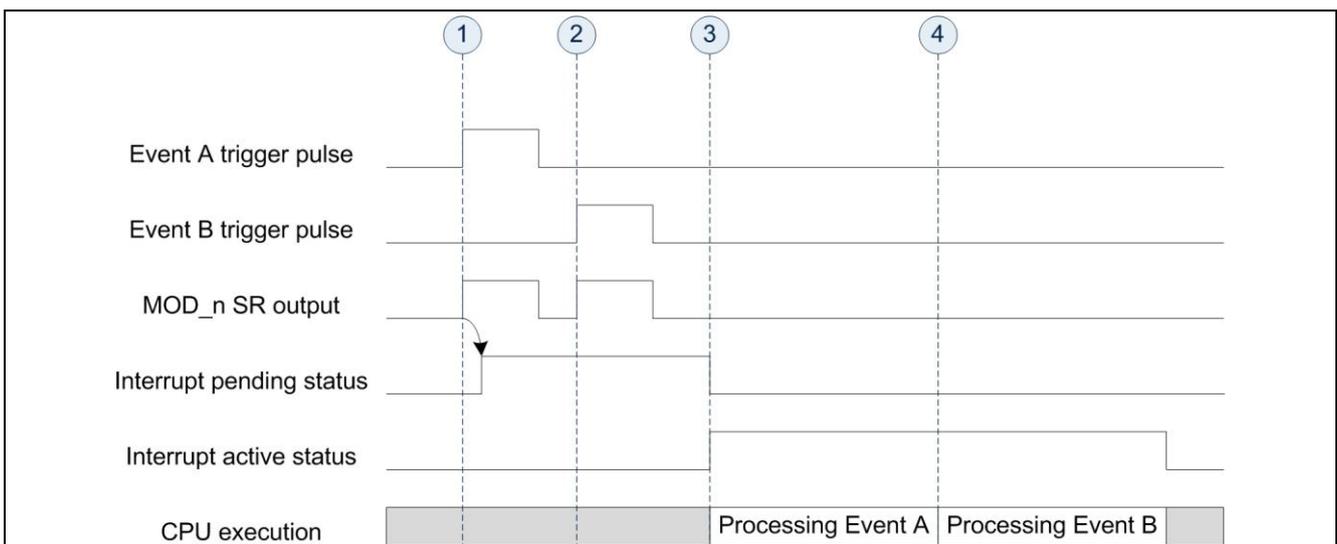
To illustrate how such an interrupt is processed, consider the following three scenarios:

- Events A and B occur before interrupt is serviced
- Event B occurs while servicing interrupt event A
- Event A occurs while servicing interrupt event B

### Events A and B occur before interrupt is serviced

Figure 1 shows the case where events A and B occur before the interrupt is serviced. The time instances 1 to 4 are:

1. Event A generates an interrupt pulse at the module service request output and causes the interrupt to switch to a 'pending' state.
2. Event B similarly generates an interrupt pulse. However, as the interrupt is already pending due to event A, the second interrupt pulse has no additional effect.
3. CPU enters the interrupt handler, sees that event A status flag is set and executes the corresponding handler code. The interrupt changes from pending to active state.
4. The CPU checks for event B status flag and sees that it is also set. Therefore, the event B handler code will also be executed within the same instance of interrupt entry.

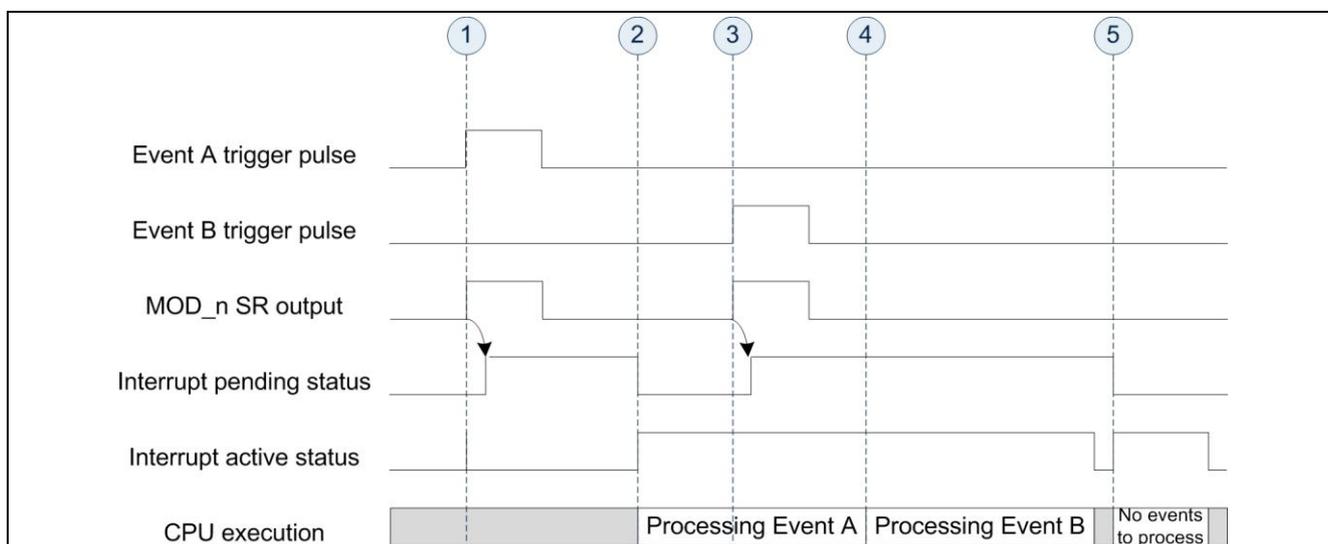


**Figure 1 Events A and B occur before interrupt is serviced**

### Event B occurs while servicing interrupt event A

Figure 2 shows the case where event B becomes pending while the CPU is still servicing the interrupt triggered by event A. The time instances 1 to 5 are:

1. Event A generates an interrupt pulse at the module service request output and causes the interrupt to switch to a 'pending' state.
2. CPU enters the interrupt handler, sees that event A status flag is set and executes the corresponding handler code. The interrupt changes from pending to active state.
3. Event B generates an interrupt pulse that changes the interrupt state to be active and pending.
4. The CPU checks for event B status flag and sees that it is also set. Therefore, the event B handler code will also be executed within the same instance of interrupt entry. Once this is done, the interrupt active state is removed.
5. However, as there is still an interrupt pending (due to event B), the CPU enters the interrupt handler again. Since all event status flags have already been cleared, the handler exits without processing any events; i.e. a dummy interrupt.

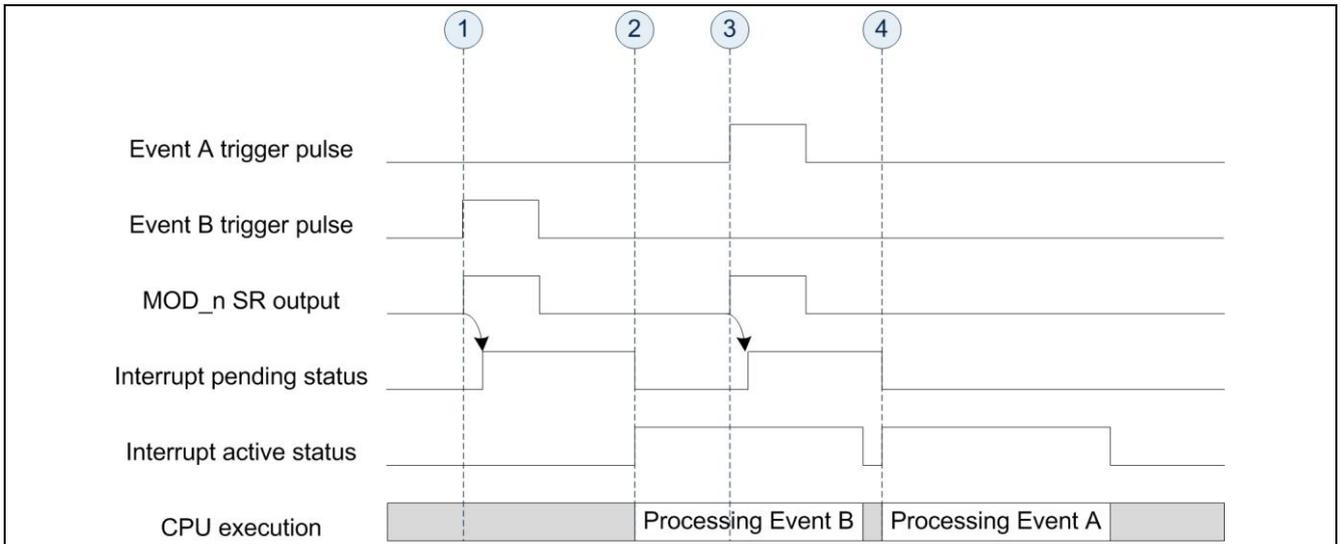


**Figure 2 Event B becomes pending while servicing interrupt A**

### Event A occurs while servicing interrupt event B

Figure 3 shows the case where event A becomes pending while the CPU is still servicing the interrupt triggered by event B. The time instances 1 to 4 are:

1. Event B generates an interrupt pulse at the module service request output and causes the interrupt to switch to a 'pending' state.
2. CPU enters the interrupt handler and causes the interrupt to change from pending to active state. CPU detects that only event B status flag is set and executes the corresponding handler code.
3. Event A generates an interrupt pulse that changes the interrupt state to be active and pending.
4. After executing the handler code for event B, the CPU exits the interrupt handler. However, as there is still an interrupt pending (due to event A), the CPU enters the interrupt handler again. The CPU detects that event A status flag is set and executes the corresponding handler code.



**Figure 3 Event A becomes pending while servicing interrupt B**

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